



University
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Regret-equality in Stable Marriage

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Outline

- Matching problems
- Fairness
- Finding fair stable matchings
- Experiments
- Future work

Matching Problems



- Assign one set of entities to another set of entities
- Based on preferences and capacities

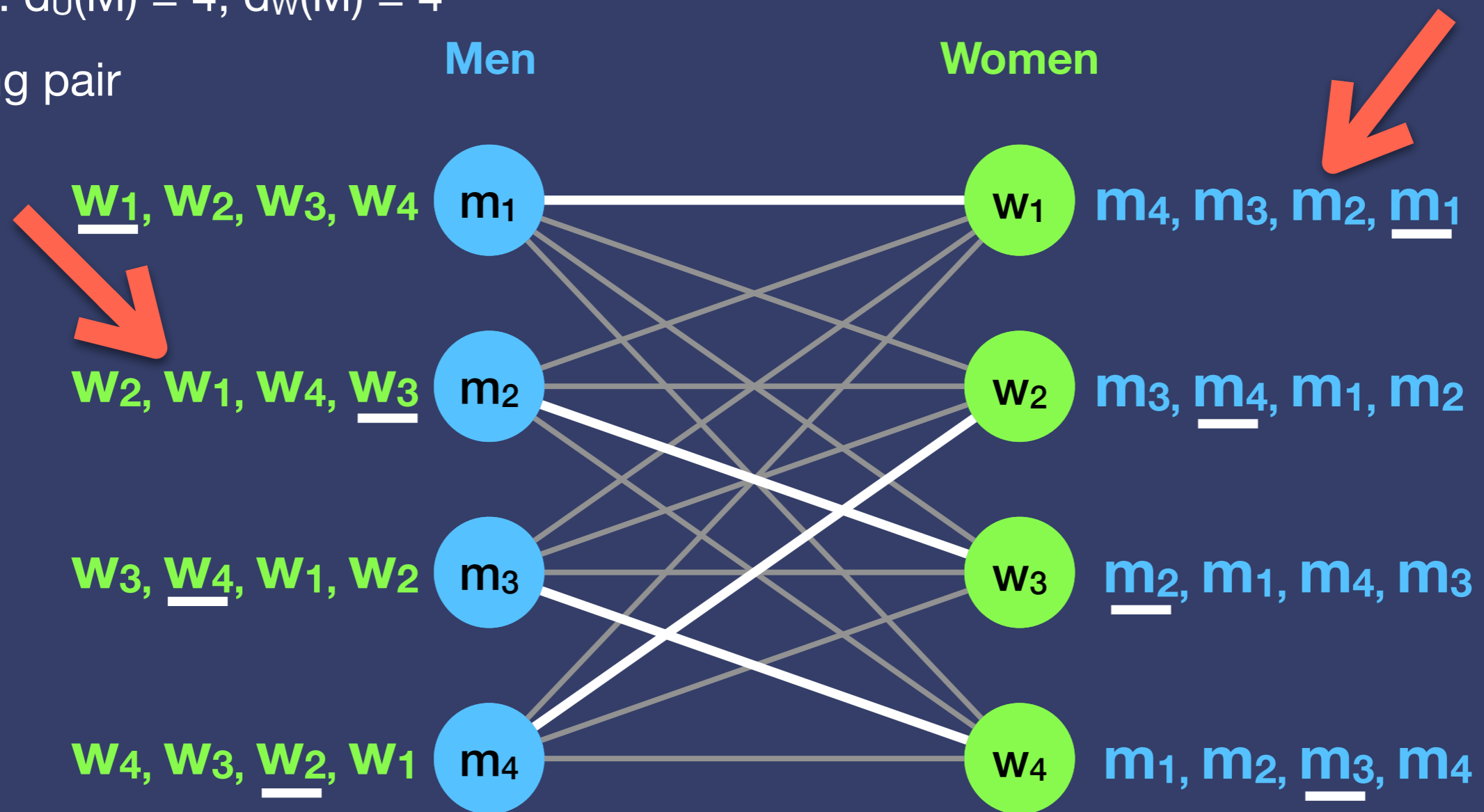
Rank

Stable Marriage

Cost: $c_U(M) = 10$, $c_W(M) = 10$

Degree: $d_U(M) = 4$, $d_W(M) = 4$

Blocking pair



A **stable matching** is a matching with no blocking pairs

Stable Marriage

- A **stable matching** is a matching with no blocking pairs
- Many stable matchings per instance
- We can find a stable matching in linear time using the man-oriented or woman-oriented Gale-Shapley Algorithm. $O(m)$ time where m is total length of preference lists
- Man-oriented Gale-Shapley Algorithm: finds a man-optimal (woman-pessimal) stable matching (and vice versa)

Fairness

- Want to find a stable matching that provides some kind of equality between men and women
- Several different fairness measures



Fairness measures

Among all stable matchings, find the stable matching that...

Cost: $c_U(M)$, $c_W(M)$

Degree: $d_U(M)$, $d_W(M)$

Minimises the maximum

balanced score

Balanced stable matching NP-hard

degree

Minimum-regret stable matching Poly

Minimises the difference

sex-equal score

Sex-equal stable matching NP-hard

regret-equal score

* Regret-equal stable matching ?

Minimises the sum

egalitarian cost

Egalitarian stable matching Poly

regret sum score

* Min-regret sum stable matching ?

Fairness measures (degree based)

10 stable matchings for this instance

m_1 : $w_1, w_2, \underline{w_3}, w_4$	w_1 : $m_4, m_3, \underline{m_2}, m_1$
m_2 : $w_2, \underline{w_1}, w_4, w_3$	w_2 : $m_3, \underline{m_4}, m_2, m_1$
m_3 : $w_3, \underline{w_4}, w_1, w_2$	w_3 : $m_2, \underline{m_1}, m_4, m_3$
m_4 : $w_4, w_3, \underline{w_2}, w_1$	w_4 : $m_1, m_2, \underline{m_3}, m_4$

m_1 : $w_1, w_2, \underline{w_3}, w_4$	w_1 : $m_4, \underline{m_3}, m_2, m_1$
m_2 : $w_2, w_1, \underline{w_4}, w_3$	w_2 : $m_3, \underline{m_4}, m_2, m_1$
m_3 : $w_3, w_4, \underline{w_1}, w_2$	w_3 : $m_2, \underline{m_1}, m_4, m_3$
m_4 : $w_4, w_3, \underline{w_2}, w_1$	w_4 : $m_1, \underline{m_2}, m_3, m_4$

m_1 : $\underline{w_1}, w_2, w_3, w_4$	w_1 : $m_4, m_3, m_2, \underline{m_1}$
m_2 : $\underline{w_2}, w_1, w_4, w_3$	w_2 : $m_3, m_4, \underline{m_2}, m_1$
m_3 : $\underline{w_3}, w_4, w_1, w_2$	w_3 : $m_2, m_1, m_4, \underline{m_3}$
m_4 : $\underline{w_4}, w_3, w_2, w_1$	w_4 : $m_1, m_2, m_3, \underline{m_4}$

Min-regret & Regret-equal

Degree: 3

Regret-equality score: 0

Min-regret sum score: 6

Min-regret & Min-regret sum

Degree: 3

Regret-equality score: 1

Min-regret sum score: 5

Min-regret sum

Degree: 4

Regret-equality score: 3

Min-regret sum score: 5

Over all stable matchings:

Minimum degree = 3

Minimum regret-equality score = 0

Minimum regret sum score = 5

Finding a Regret-Equal Stable Matching



Rotations

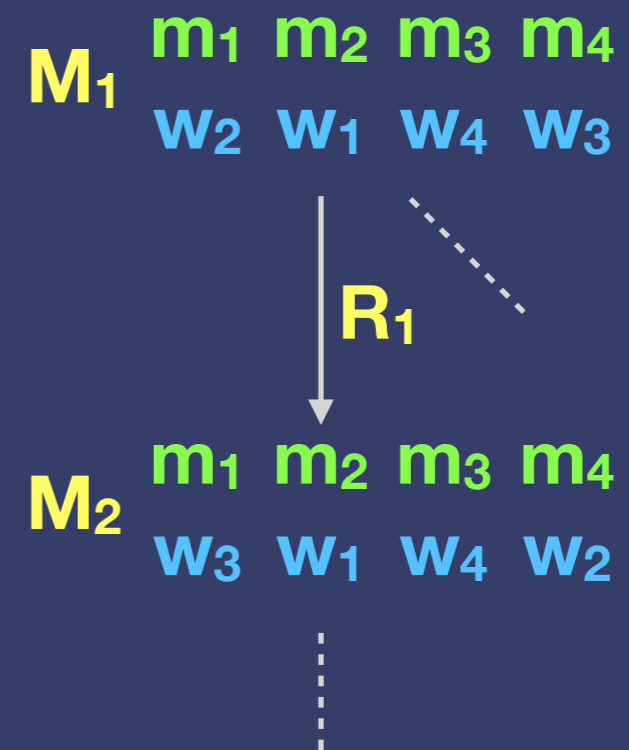
- Rotation - series of man-woman pairs that take us from one stable matching to another when permuted

R_1 m_1 m_4
 w_2 w_3

- Can only eliminate *exposed* rotations

R_2 m_1 m_2
 w_1 w_2

- $O(n^2)$ algorithm to find all rotations
- Rotations form a structure to allow enumeration of all stable matchings. All rotation makes some men worse off and some women better off



Algorithm

1. Find the man-optimal stable matching M_0

- Each man has their best partner in any stable matching.

Say $d_u(M_0) = 2$ and $d_w(M_0) = 5$ $d(M_0) = (2, 5)$

- Then, a regret equal stable matching must exist within the following degrees pairs:

r-e score: 3 $(2, 5)$

r-e score: 2 $(2, 4) (3, 5)$

r-e score: 1 $(2, 3) (3, 4) (4, 5)$

r-e score: 0 $(2, 2) (3, 3) (4, 4) (5, 5)$

r-e score: 1 $(2, 1) (3, 2) (4, 3) (5, 4) (6, 5)$

r-e score: 2 $(3, 1) (4, 2) (5, 3) (6, 4) (7, 5)$

why are these the only possible degrees?

- M_0 has a r-e score of 3
- men can only get worse
- women can only get better

Algorithm

2. If $d_U(M_0) \geq d_W(M_0)$ then exit with M_0
3. For each man m and for each column c :
 1. rotate m down to c (if possible)
 2. rotate women down column c who have worst rank
 - Stop iterating women up the column when $d_U(M) \geq d_W(M)$
 - Save the best matching as you go

r-e score: 3 (2, 5)

r-e score: 2 (2, 4) (3, 5)

r-e score: 1 (2, 3) (3, 4) (4, 5)

r-e score: 0 (2, 2) (3, 3) (4, 4) (5, 5)

r-e score: 1 (2, 1) (3, 2) (4, 3) (5, 4) (6, 5)

r-e score: 2 (3, 1) (4, 2) (5, 3) (6, 4) (7, 5)

Time complexity

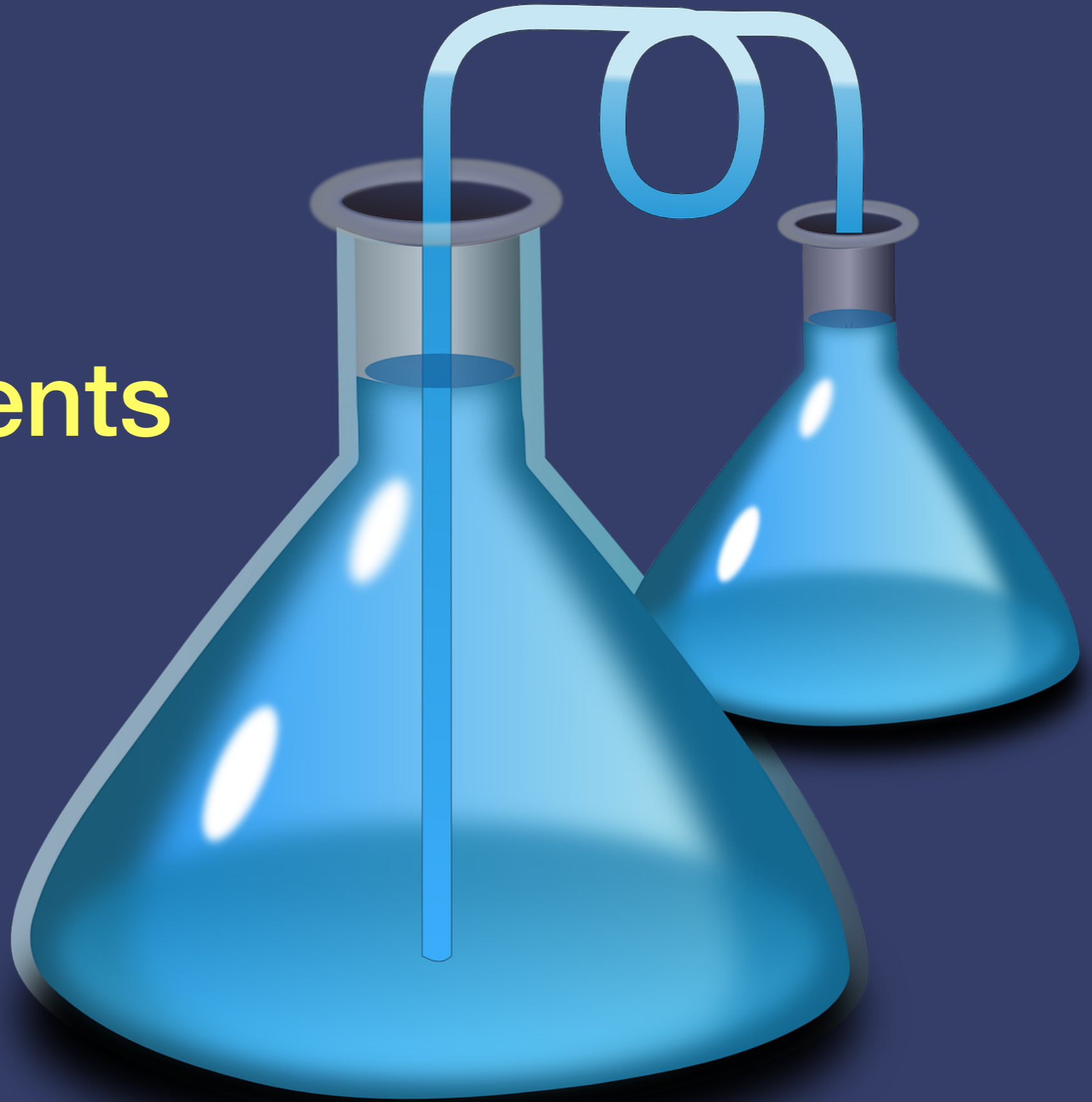
- Find man-optimal stable matching & all rotations $O(n^2)$
- For each man $O(n)$
 - For each column $O(2 * \text{man-optimal difference} * |d_u(M_0) - d_w(M_0)|) = O(c)$
 - Rotate man down and women down $O(n^2)$



Total $O(n^3c)$



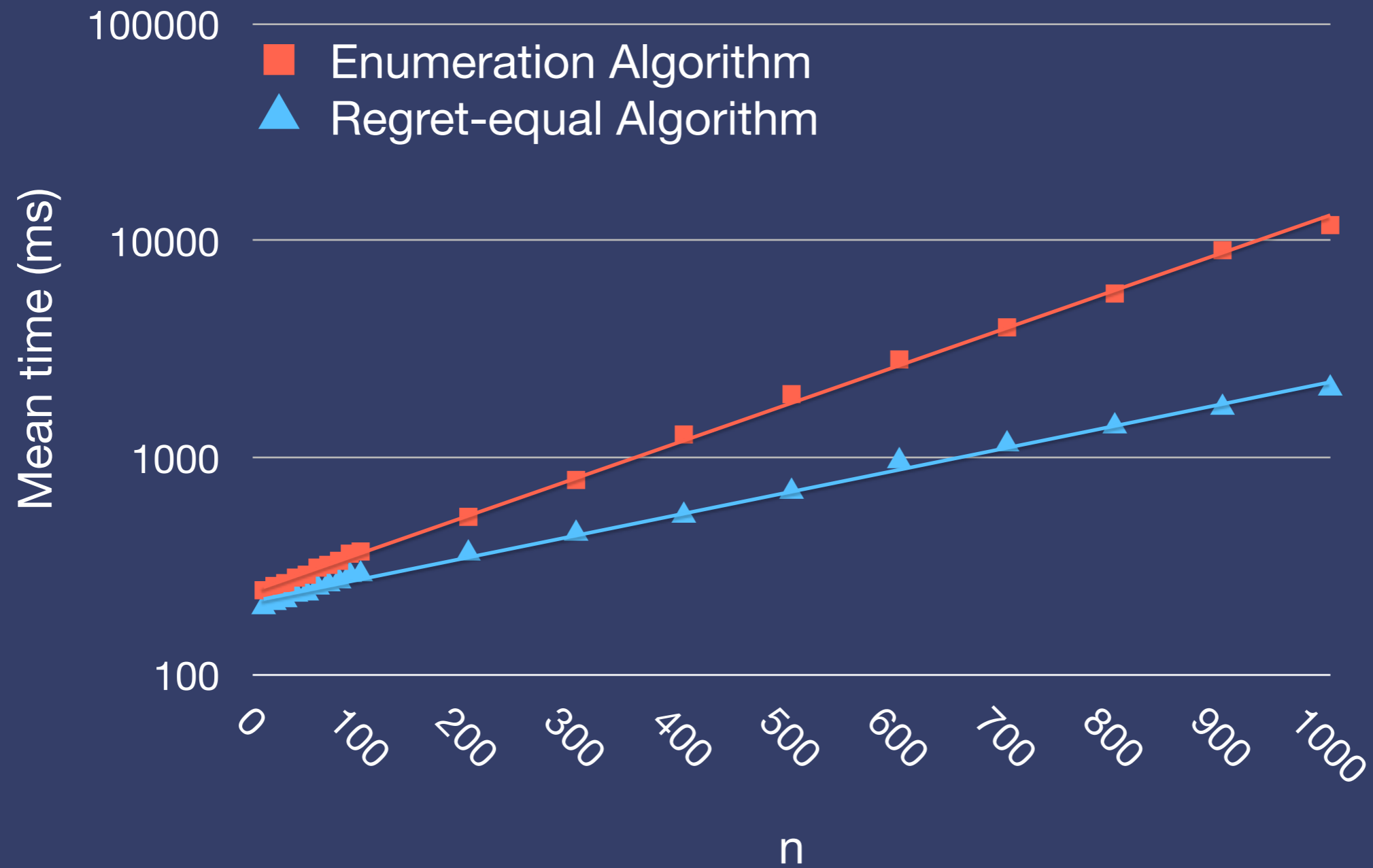
Experiments



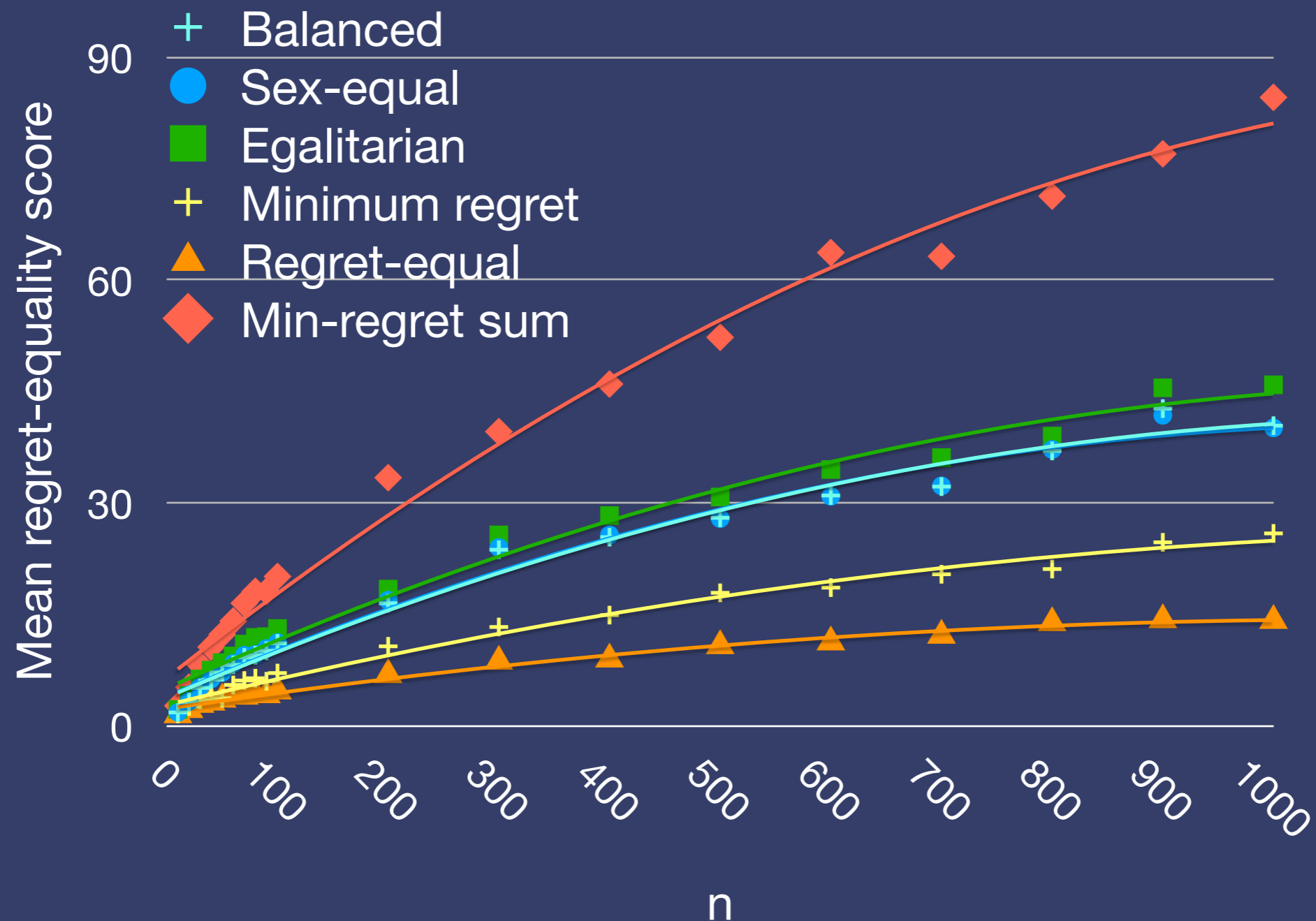
Methodology

- Performance of the Regret-equal Algorithm compared to an Enumeration algorithm (exponential in worst case)
- Instances size {10, 20, ..., 100, 200, ..., 1000}, complete preference lists, 500 instance per size.
- looked at properties over several types of optimal stable matching (balanced, sex-equal, egalitarian, minimum regret, regret-equal, min-regret sum)
- Java, Python, Bash, GNU parallel
- Correctness
 - all matchings found were stable
 - Regret-equality scores matched
 - CPLEX up to size $n = 50$ for the enumeration algorithm

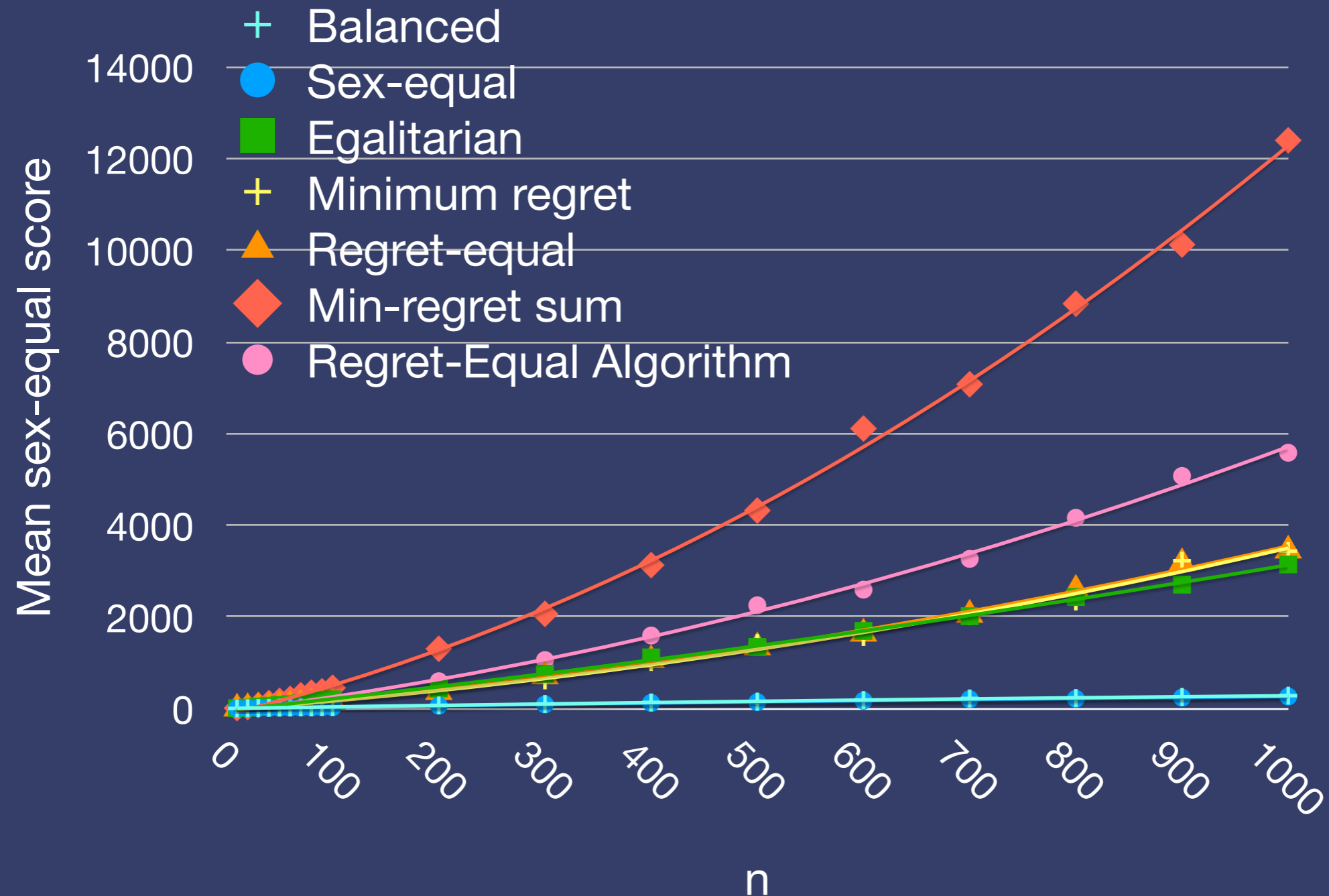
Time taken



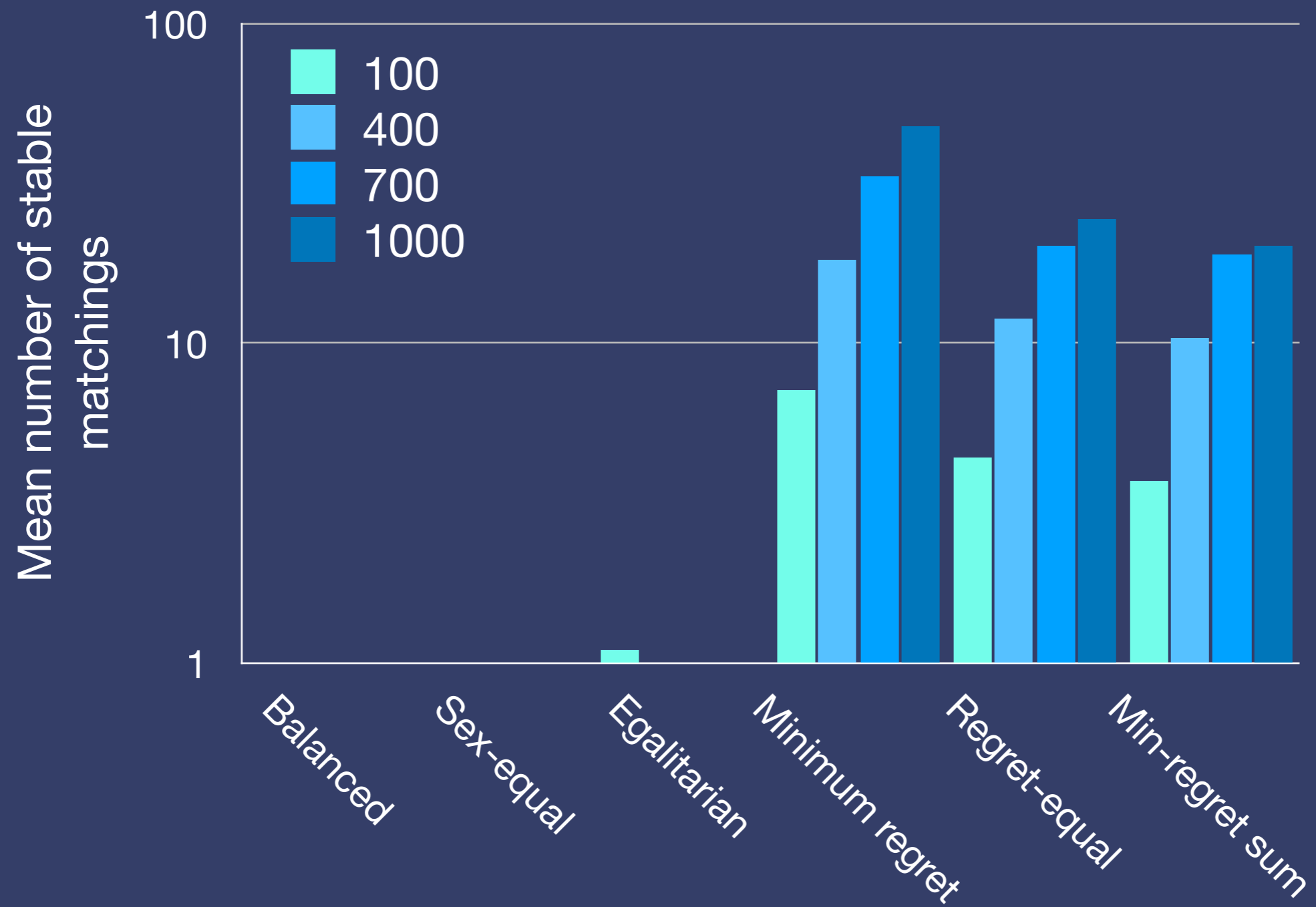
Regret-equality score for different optimal matchings



Sex-equal score for different optimal matchings



Frequency of different optimal stable matchings



Future Work

- Improving the $O(n^3c)$ Regret-equal Algorithm, where $c = |d_U(M_0) - d_W(M_0)|$
- Grouping women - e.g. women are workers and men are jobs to assign to workers.
 - Woman optimal stable matching would naturally satisfy ‘balanced’, ‘min-regret’, ‘egalitarian’ and ‘min-regret sum’ criteria
 - Can find a ‘regret-equal’ stable matching in $O(n^4)$ time
 - Open problem for ‘sex-equality’ \rightarrow grouped-women-equality

Thank you

Summary

- Matching problems
- Fairness
- Finding fair stable matchings
- Experiments
- Future work: finding improved algorithms



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